

# An Introduction to Higher Categories

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**Bibliography: J. Lurie's HTT & HA**

Ideally, an  $\infty$ -category is a mathematical object that assembles:

- *objects*,
- *1-morphisms between objects*,
- for every  $n \geq 2$ , a notion of *n-morphisms* between  $n - 1$ -morphisms;
- for every  $n \geq 1$ , (weak) composition laws of  $n$ -morphisms only well-defined up to the data of higher morphisms.
- Associativity of compositions up to homotopy:

**Proto-Example** (Fundamental  $\infty$ -groupoid)  $X$  a CW-complex.

- Objects = points,
- 1-morphisms=paths,
- 2-morphisms=homotopy of paths ( 2-cells);
- 3-morphisms= homotopies between homotopies of paths (3-cells)
- ...

# Models:

**Problem:** No direct definition of higher categories simultaneously operational and close to this platonic form (infinite axioms!).

**First Breakthrough :** Avoid the problem by modeling the platonic form of  $\infty$ -categories using "exaggerated" templates/models that contain more structure than what the platonic form requires.

**Formally:** Find a model category whose objects serve as models for  $\infty$ -categories (Quasi-categories, Segal Categories, Simplicial Categories, etc).

**Modeling is a common practice:**

- Homotopy Theory of Spaces (Homotopy Types)  $\rightsquigarrow$  Modeled by topological spaces, simplicial sets, categories, etc
- Homotopy Theory of homotopy-commutative algebras over  $\mathbb{Q}$ :  $\rightsquigarrow$  Modeled by simplicial algebras, diff. graded algebras.
- Derived and Higher Stacks (Modeled by simplicial presheaves);

# Models: why important?

**Question:** What is the fundamental role of models and why are there are so many for the same theory?

**Second Breakthrough:** Every model category has an associated  $\infty$ -category which captures all the important information of the model structure ([Dwyer-Kan Localization](#))

**Models play a double role:** (operational) Need ambient model to shape  $\infty$ -categories; (fundamental) Every other model incarnates as an object of this ambient model.

**Consequence:** Plenty of examples of wannabe  $\infty$ -categories:

- $\exists$   $\infty$ -category of [spaces](#)  $\mathcal{S}$ : Model structure on topological spaces/simplicial sets;
- $\infty$ -category of [derived affine schemes](#) and [derived stacks](#);
- $\infty$ -category of [chain complexes up to quasi-iso.](#)

**Question:** If we already have the explicit models, why do we care about their associated  $\infty$ -categories?

**Answer:**

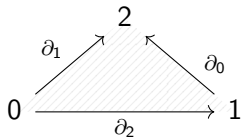
- Not all  $\infty$ -categories have a practical model presentation (Typical Examples:  $\infty$ -categories of algebra-objects in a  $\otimes$  –  $\infty$ -category: ring spectra);
- There is no sufficiently refined notion of functor to relate different models. The relevant notion is that of  $\infty$ -functor between the  $\infty$ -categories associated to the models.
- Models for diagrams are not in general given by diagrams of models.

# Modeling with Simplicial Sets: Notations

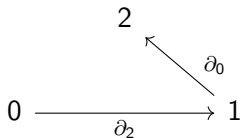
Category  $\Delta$ : objects = finite ordered sets  $[n] = \{0 < 1 < \dots < n\}$ ;  
 morphisms = order-preserving maps;  $S\text{Sets} := \text{Fun}(\Delta^{\text{op}}, \text{Sets})$ .

$$\begin{array}{ccccccc}
 & \xrightarrow{\partial_3} & & \xrightarrow{\partial_2} & & \xrightarrow{\partial_1} & \\
 \xleftarrow{\epsilon_2} & & \xleftarrow{\epsilon_1} & & \xleftarrow{\epsilon_0} & & \\
 \xrightarrow{\partial_2} & & \xrightarrow{\partial_1} & & \xrightarrow{\partial_0} & & \\
 \xleftarrow{\epsilon_1} & & \xleftarrow{\epsilon_0} & & & & \\
 \xrightarrow{\partial_1} & & \xrightarrow{\partial_0} & & & & \\
 \xleftarrow{\epsilon_0} & & & & & & \\
 \xrightarrow{\partial_0} & & & & & & \\
 \dots & S_3 & & S_2 & & S_1 & & S_0
 \end{array}$$

$\Delta[n] := \text{Yoneda}([n])$ ;  $S_n = \text{Hom}(\Delta[n], S)$ ;  $\Delta[2] :=$



$\Lambda_n^j := \Delta[n]$  - the  $j$ th face and the interior ;  $\Lambda_2^1 =$



$S$  simplicial set  $\rightsquigarrow |S|$  its topological realization; Equivalences of simplicial sets := weak-homotopy equivalences ( $\pi_\bullet(|-|)$ -eq.)

# Modeling with Simplicial Sets

**Modeling usual categories:**  $\mathcal{C}$  usual category  $\rightsquigarrow$  simplicial set  $N(\mathcal{C})$  (Nerve) with

$$\{n\text{-simplexes}\} := \{\text{strings } X_0 \xrightarrow{f_0} X_1 \xrightarrow{f_1} \dots \xrightarrow{f_{n-1}} X_n\}$$

Boundary maps  $\partial_i$  encode composition law. Degeneracy maps  $\epsilon_i$  encode identity maps.

- $\{\text{Functors } \mathcal{C} \rightarrow \mathcal{D}\} \simeq \{\text{Simplicial Maps } N(\mathcal{C}) \rightarrow N(\mathcal{D})\}$
- $X \simeq N(\mathcal{C}) \leftrightarrow \forall n \geq 2, \forall 0 < i < n, \forall u : \Lambda_n^i \rightarrow X$

$$\begin{array}{ccc} \Lambda_n^i & \xrightarrow{u} & X \\ \downarrow & \nearrow & \\ \Delta[n] & & \end{array} \quad \exists! \text{ factorization}$$

# Modeling with Simplicial Sets

**Generating compositions:** Take diagram shapes

$$\Lambda_2^1 = \begin{array}{ccc} & 2 & \\ & \swarrow \partial_0 & \\ 0 & \xrightarrow{\partial_2} & 1 \end{array} \subseteq \Delta[2] := \begin{array}{ccc} & 2 & \\ \nearrow \partial_1 & & \swarrow \partial_0 \\ 0 & \xrightarrow{\partial_2} & 1 \end{array}$$

$\Lambda_2^1 \rightarrow N(\mathcal{C}) \leftrightarrow$  string of morphisms  $X \xrightarrow{f} Y \xrightarrow{g} Z$  in  $\mathcal{C}$

$$\begin{array}{ccc} \Lambda_2^1 & \longrightarrow & N(\mathcal{C}) \\ \downarrow \wr & \nearrow & \uparrow \\ \Delta[2] & & \end{array} \quad \begin{array}{c} \leftrightarrow \exists! g \circ f := \partial_1 \\ \begin{array}{ccc} & Z & \\ \nearrow & & \swarrow g \\ X & \xrightarrow{f} & Y \end{array} \end{array}$$

$\exists!$  extension

$\exists!$  extension  $\leftrightarrow \exists!$  compositions  $\leftrightarrow$  shaded faces



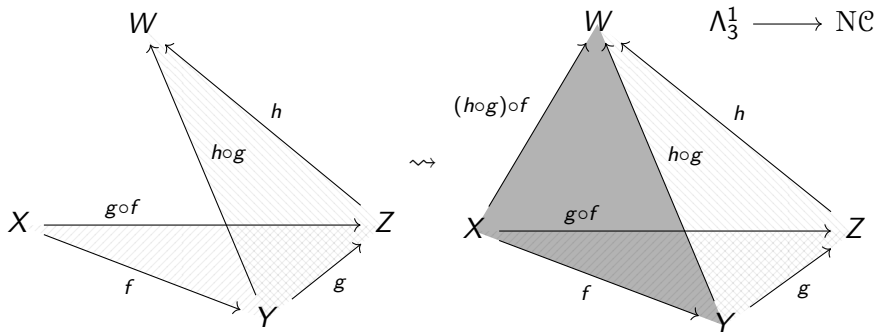
# Modeling with Simplicial Sets

## Processing Associativity:

$$X \xrightarrow{f} Y \xrightarrow{g} Z \xrightarrow{h} W \text{ in } \mathcal{C} \leftrightarrow \Lambda_2^1 \xrightarrow{(g,f)} N(\mathcal{C}) + \Lambda_2^1 \xrightarrow{(h,g)} N(\mathcal{C})$$

Generate Compositions:  $\Delta[2] \xrightarrow{(g,f)} N(\mathcal{C}) + \Delta[2] \xrightarrow{(h,g)} N(\mathcal{C})$

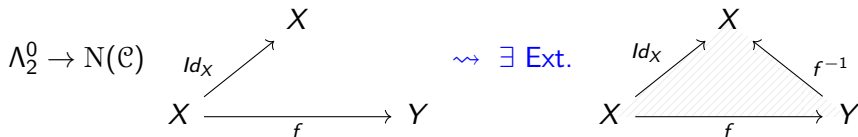
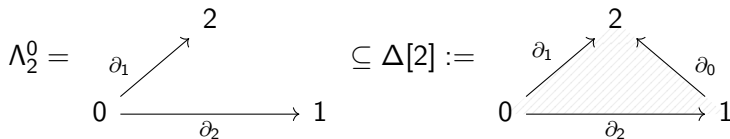
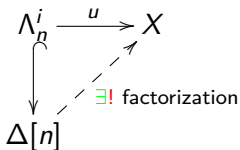
Glue



$$(h \circ g) \circ f = h \circ (g \circ f) \leftrightarrow \text{Filling back face} \leftrightarrow \text{Extend along } \Lambda_3^1 \subseteq \Delta[3]$$

# Modeling with Simplicial Sets

**Groupoids:**  $X \simeq N(\mathcal{C} \text{ groupoid}) \leftrightarrow \forall n \geq 2, \forall 0 \leq i \leq n, \forall u : \Lambda_n^i \rightarrow X$

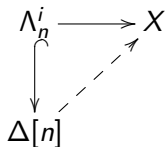


**Summary:** Allowing lifting property for extremes  $i = 0, n$  gives inverses. For terms in the middle  $0 < i < n$  gives compositions.

# Modeling with Simplicial Sets

**Modeling spaces:**  $T$  topological space  $\rightsquigarrow$  simplicial set  $\text{Sing}(T)$  with  $n$ -simplexes given by continuous maps from the topological  $n$ -simplex  $\Delta^n$  to  $T$

$X \simeq \text{Sing}(T) \rightarrow \forall n \geq 2, \forall 0 \leq i \leq n, \forall u : \Lambda_n^i \rightarrow X \exists$  factorization not necessarily unique:



**Slogan:** Kan-complexes are those simplicial sets where the direction of the arrows is irrelevant.

# Modeling with Simplicial Sets

**Modeling  $\infty$ -categories:** A Quasi-category is a simplicial set which shares simultaneously features of categories and spaces:

## Definition

A Quasi-category is a simplicial set  $\mathcal{C}$  with the following property:

$$\forall n \geq 2, \forall 0 < i < n, \forall u : \Lambda_n^i \rightarrow \mathcal{C}$$

$$\begin{array}{ccc} \Lambda_n^i & \longrightarrow & \mathcal{C} \\ \downarrow & \nearrow \text{---} & \nearrow \\ \Delta[n] & & \end{array} \quad \exists \text{ factorization (not nec. unique!)}$$

## Examples:

- $\mathcal{C}$  category  $\implies$   $N(\mathcal{C})$  quasi-category with uniquely-defined compositions;
- $X$  Kan complex  $\implies$   $X$  quasi-category;

# Playing with Quasi-categories: Unveiling the definition

$\mathcal{C}$  quasi-category. How far can we go?

**Generating "Compositions":**

$$\Lambda_2^1 = \begin{array}{ccc} & 2 & \\ & \swarrow \partial_0 & \\ 0 & \xrightarrow{\partial_2} & 1 \end{array} \subseteq \Delta[2] := \begin{array}{ccc} & 2 & \\ \nearrow \partial_1 & & \nwarrow \partial_0 \\ 0 & \xrightarrow{\partial_2} & 1 \end{array}$$

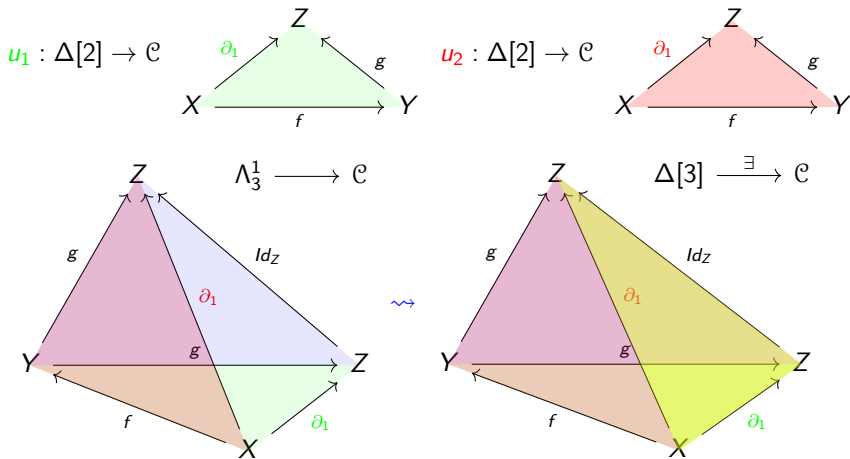
$\Lambda_2^1 \rightarrow \mathcal{C} \leftrightarrow$  string of 1-simplices  $X \xrightarrow{f} Y \xrightarrow{g} Z$  in  $\mathcal{C}$

$$\begin{array}{ccc} \Lambda_2^1 & \longrightarrow & \mathcal{C} \\ \downarrow & \nearrow & \uparrow \\ \Delta[2] & \xrightarrow{\text{extension}} & \mathcal{C} \end{array} \quad \leftrightarrow \quad \begin{array}{ccc} & Z & \\ \nearrow \partial_1 & & \nwarrow g \\ X & \xrightarrow{f} & Y \end{array}$$

Each 2-simplex in  $\mathcal{C}$  "makes a form of commutativity". There can be many!

# Playing with Quasi-categories: Unveiling the definition

**Control of non-uniqueness of "compositions":**



**2-simplex**  $\leftrightarrow$  homotopy of compositions  $\partial_1 \sim \partial_1$ . Lifting Mechanism provides compatibility between compositions.

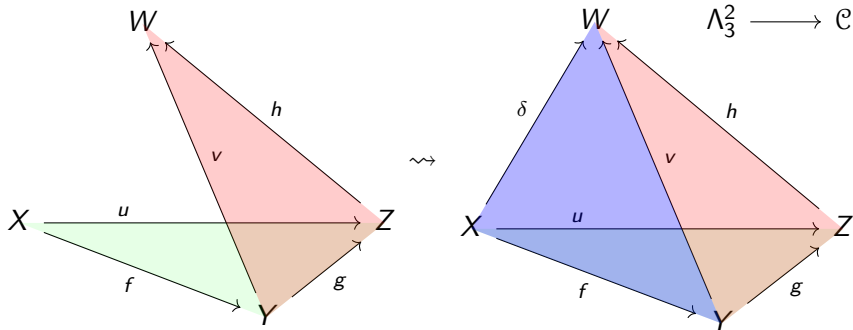
**More Generally:** compatibility of  $n$  different compositions  $\rightsquigarrow$   $n$ -simplexes.

# Playing with Quasi-categories: Unveiling the definition

## Processing "Associativity":

$$X \xrightarrow{f} Y \xrightarrow{g} Z \xrightarrow{h} W \text{ in } \mathcal{C} \leftrightarrow \Lambda_2^1 \xrightarrow{(g,f)} \mathcal{C} + \Lambda_2^1 \xrightarrow{(h,g)} \mathcal{C}$$

Exhibit compositions:  $\Delta[2] \xrightarrow{(g,f)} \mathcal{C} + \Delta[2] \xrightarrow{(h,g)} \mathcal{C}$



$$\delta \sim "h \circ u" \leftrightarrow \text{Filling back face} \leftrightarrow \text{Extend along } \Lambda_3^1 \subseteq \Delta[3]$$

# Quasi-categories as $\infty$ -categories: $\infty$ -functors

**Definition:** An  $\infty$ -category is a Quasi-category. **Objects** = 0-simplexes. **Morphisms** = 1-simplexes.

**Definition:** An  $\infty$ -functor is a map of simplicial sets between quasi-categories.

**Explanation:** Functors  $\leftrightarrow$  functions that preserve commutative diagrams  $\leftrightarrow$  send  $n$ -simplexes to  $n$ -simplexes and preserve boundaries.

**Prop. ( $\infty$ -category of functors)**  $\mathcal{C}$   $\infty$ -category,  $K$  simplicial set (diagram shape):

$$\mathrm{Fun}(K, \mathcal{C}) := \underline{\mathrm{Hom}}_{\Delta}(K, \mathcal{C}) \quad (\text{internal-hom})$$

is a  $\infty$ -category.

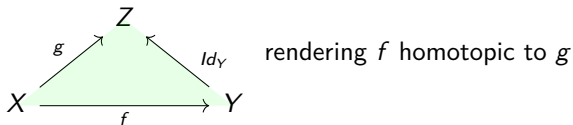
**Prop. (Products)**  $\mathcal{C}, \mathcal{D}$   $\infty$ -categories  $\implies \mathcal{C} \times \mathcal{D}$   $\infty$ -category.



# Quasi-categories as $\infty$ -categories: Homotopy Category

$\mathcal{C}$   $\infty$ -category  $\rightsquigarrow$  truncation (forgets higher cells) produces usual category  $\mathbf{h}\mathcal{C}$  (homotopy category):

- Objects of  $\mathbf{h}\mathcal{C}$ : 0-simplexes of  $\mathcal{C}$ ;
- Morphisms of  $\mathbf{h}\mathcal{C}$ : := homotopy classes of 1-morphisms:  $f, g$  are **equivalent** iff there exists a 2-morphism  $u : \Delta[2] \rightarrow \mathcal{C}$



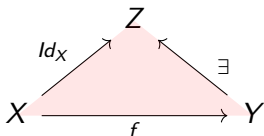
- Compositions: Well-defined using the lifting property.

**Definition:** **Subcategory** of a quasi-category  $\mathcal{C}$  is a sub-simplicial set  $\mathcal{C}'$  obtained as a fiber product in simplicial sets

$$\begin{array}{ccc} \mathcal{C}' & \longrightarrow & \mathcal{C} \\ \downarrow & & \downarrow \\ \mathbf{N}(D) & \longrightarrow & \mathbf{N}(\mathbf{h}\mathcal{C}) \end{array} \quad \text{with } D \text{ a subcategory of } \mathbf{h}\mathcal{C}.$$

# Quasi-categories as $\infty$ -categories: Equivalences

**Definition:** A 1-morphisms  $f$  of  $\mathcal{C}$  is said to be an **equivalence** if its homotopy class  $[f]$  in  $\mathbf{h}\mathcal{C}$  is an isomorphism, ie,  $\exists$



rendering  $f$  invertible up to a 2-cell

**Definition:** An  $\infty$ -category is an  **$\infty$ -groupoid** if all its 1-morphisms are invertible.

**Prop:** An  $\mathcal{C}$   $\infty$ -category is an  $\infty$ -groupoid if and only if it is a Kan-complex.

**Example:**  $T$  topological Space  $\rightsquigarrow$  (Proto-Example) Fundamental  $\infty$ -groupoid of  $T := \text{Sing}(T)$ .

# Quasi-categories as $\infty$ -categories: Mapping Spaces

$\mathcal{C}$   $\infty$ -category.  $X, Y$  objects  $\rightsquigarrow \exists$  "space" of morphisms  $X \rightarrow Y$ ,

**Definition:** The **mapping space** between  $X$  and  $Y$  is the simplicial set  $\text{Map}_{\mathcal{C}}(X, Y)$  defined as the fiber product

$$\begin{array}{ccccc}
 \Delta[0] & & \text{Fun}(\Delta[1], \mathcal{C}) & & \Delta[0] \\
 \searrow^X & & \swarrow^{ev_0} & & \swarrow^Y \\
 & & \mathcal{C} & & \mathcal{C} \\
 & & \swarrow^{ev_1} & & \searrow^Y
 \end{array}$$

**Proposition:** Let  $\mathcal{C}$  be an  $\infty$ -category. Then  $\text{Map}_{\mathcal{C}}(X, Y)$  is a Kan-complex. ( $\Leftrightarrow$  "for  $n \geq 2$ ,  $n$ -morphisms are invertible").

0-simpl. :=  $f : \Delta[1] \rightarrow \mathcal{C}$  with  $\partial_1(f) = X$ ,  $\partial_0(f) = Y$ .

1-simpl. :=  $\Delta[1] \times \Delta[1] \rightarrow \mathcal{C}$

$$\begin{array}{ccc}
 X & \xrightarrow{f} & Y \\
 \downarrow & \searrow & \downarrow \\
 X & \xrightarrow{f} & Y
 \end{array}$$

,  $n$ -simpl. :=  $\Delta[n] \times \Delta[1] \rightarrow \mathcal{C}$

**Remark:**  $\text{Hom}_{\text{hc}}(X, Y) = \pi_0(\text{Map}_{\mathcal{C}}(X, Y))$

# Quasi-categories as $\infty$ -categories: Mapping Spaces & Compositions

**Warning:** There is no **strict** manifestation of the "composition law" in  $\mathcal{C}$ ,  $\text{Map}_{\mathcal{C}}(X, Y) \times \text{Map}_{\mathcal{C}}(Y, Z) \not\rightarrow \text{Map}_{\mathcal{C}}(X, Z)$

**Rectification (Lurie)**  $\exists$  alternative simplicial sets  $\widetilde{\text{Map}}_{\mathcal{C}}(X, Y)$ , together with a canonical zig-zags of weak-equivalences of SSets

$$\widetilde{\text{Map}}_{\mathcal{C}}(X, Y) \xrightarrow{\sim} \bullet \xleftarrow{\sim} \text{Map}_{\mathcal{C}}(X, Y)$$

such that the "composition law" of  $\mathcal{C}$  is translated by concrete strict maps

$$\widetilde{\text{Map}}_{\mathcal{C}}(X, Y) \times \widetilde{\text{Map}}_{\mathcal{C}}(Y, Z) \rightarrow \widetilde{\text{Map}}_{\mathcal{C}}(X, Z)$$

**Good**  $\mathcal{C}$  is converted into a concrete **simplicial category** (notation:  $\mathcal{C}[\mathcal{C}]$ ) with strictly associative compositions;

**Bad**  $\widetilde{\text{Map}}_{\mathcal{C}}(X, Y)$  are not really "spaces", ie, not Kan-complexes. Very complicated.

# Quasi-categories as $\infty$ -categories: Simplicial Categories

**Rectification**  $\mathcal{C} \rightsquigarrow$  Simplicial category  $\mathfrak{C}[\mathcal{C}]$

**Question:** Simplicial Categories  $\xrightarrow{?}$  Quasi-categories?

**Construction.** There exists a non-trivial extension of the nerve construction  $N$  from usual categories to simplicial enriched categories that takes into account the simplicial structure.  $N_{\Delta}(-)$  (Simplicial Nerve).

**Application:.** Model a simplicial cat.  $E$  by a simplicial set  $N_{\Delta}(E)$

**Prop.**  $E$  enriched by Kan-complexes  $\implies N_{\Delta}(E)$  is a quasi-category with

- Objects of  $N_{\Delta}(E)$  = Objects of  $E$
- Weak-equivalences of simplicial sets

$$\text{Map}_{\mathfrak{C}}(X, Y) \rightarrow E(X, Y) \text{ (enrichment)}$$

**Question:** How rich is this dictionary Simplicial categories  $\leftrightarrow$  Quasi-categories?

# Quasi-categories as $\infty$ -categories: Equivalence of Models

$\mathfrak{C}$

## Theorem (Joyal-Lurie)

*There exists a model structure on the category  $\mathbb{S}\text{Sets}$  with*

- cofibrant-fibrant objects = Quasi-categories
- weak-equivalences = essentially surjective + weak-equivalences of mapping spaces

## Theorem (Bergner)

*There exists a model structure on the category of simpl. cats.*

- fibrant objects = simpl. cats. enriched in Kan-complexes
- weak-equivalences = essentially surjective + weak-equivalences of mapping spaces

$N_{\Delta}$

## Theorem (Lurie)

$(\mathfrak{C}, N_{\Delta})$  forms a Quillen equivalence.

# Quasi-categories as $\infty$ -categories: Plug in Examples.

Use the dictionary Simplicial Cats  $\leftrightarrow$  Quasi-categories to produce examples:

## Machine to produce Examples:

$\mathcal{M}$  simplicial model category (model structure + compatible simplicial enrichment).  $\rightsquigarrow$

$E_{\mathcal{M}} := \mathcal{M}^{\circ}$  full simplicial subcategory of cofibrant-fibrant objects in  $\mathcal{M}$  is a simplicial category enriched in Kan-complexes

$\rightsquigarrow N_{\Delta}(E_{\mathcal{M}})$  is a quasi-category

.

# Quasi-categories as $\infty$ -categories: Plug in Examples.

**First Non-trivial Example:**  $\infty$ -category of spaces  $\mathcal{S} := N_{\Delta}(E_{\mathcal{M}})$ :

- $\mathcal{M} = \mathcal{S}\text{Sets}$ ; model structure to study weak-homotopy equivalences;
- Cofibrant-Fibrant objects are Kan-complexes  $X, Y$ , etc
- Simplicial structure = simplicial set of maps  $\underline{\text{Hom}}_{\Delta}(X, Y)$  (also a Kan-complex).
- Objects of  $\mathcal{S} = \text{Kan complexes}$ ;  $\text{Map}_{\mathcal{S}}(X, Y) \sim \underline{\text{Hom}}_{\Delta}(X, Y)$
- Equivalences = weak-homotopy eq. of Kan-complexes

**Example:**  $\infty$ -category of  $\infty$ -categories  $\text{Cat}_{\infty} := N_{\Delta}(E_{\mathcal{M}})$ :

- $\mathcal{M} = \mathcal{S}\text{Sets}$ ; model structure of Lurie-Joyal is **NOT** a *simplicial model category*  $\rightsquigarrow$  can be **modified** to become one (*marked simplicial sets*).
- Objects of  $\text{Cat}_{\infty} = \text{Quasi-cat.}$ ,  $\text{Map}_{\text{Cat}_{\infty}}(\mathcal{C}, \mathcal{D}) \sim \text{space of } \infty\text{-functors.}$
- Equivalences = surjective + homotopy equivalences of maps;



# Quasi-categories as $\infty$ -categories: Plug in Examples.

**Rectification of Diagrams:**  $\mathcal{M}$  a simplicial model category *combinatorial*.  $D$  be a usual category. Then there is an equivalence of quasi-categories

$$\mathrm{Fun}(N(D), N_{\Delta}(E_{\mathcal{M}})) \simeq N_{\Delta}(E_{\mathcal{M}^D})$$

where  $\mathcal{M}^D :=$  category of  $D$ -diagrams in  $\mathcal{M}$  equipped with the projective model structure to study the homotopy theory of diagrams.

**Corollary: Third Non-trivial Example:**  $D$  a usual category.  
 $\mathcal{M} = \mathbb{S}\mathrm{Sets} +$  model structure for weak-homotopy equivalences;

$$\mathrm{Fun}(N(D^{\mathrm{op}}), \mathcal{S}) \simeq \mathrm{Fun}(N(D^{\mathrm{op}}), N_{\Delta}(E_{\mathcal{M}})) \simeq N_{\Delta}(E_{\mathcal{M}^{D^{\mathrm{op}}}})$$

l.h.s =  $\mathcal{P}(N(D)) :=$  presheaves of spaces;

r.h.s = (Projective) Model structure on Simplicial Presheaves

# Quasi-categories as $\infty$ -categories: Limits and Colimits.

$\mathcal{C}$   $\infty$ -category  $\rightsquigarrow$  theory of **limits** and **colimits** internal to the language of quasi-categories.

**Definition:**  $X \in \mathcal{C}$  **initial object** iff for every object  $Z \in \mathcal{C}$  the Kan-complex  $\text{Map}_{\mathcal{C}}(X, Z)$  is **contractible**.

**Slogan:** Universal properties are defined only up to a contractible space of choices.

**Cones:**  $K$  simplicial set (diagram shape),  $d : K \rightarrow \mathcal{C}$   $\infty$ -functor (diagram). Construct a new simplicial set  $K^{\triangleright}$  by formally adding an exterior vertex to  $K$ . A "cone under  $d$ " is a map of simplicial sets  $\tilde{d} : K^{\triangleright} \rightarrow \mathcal{C}$  whose restriction to  $K$  is  $d$ .

**Prop.:** There exists a quasi-category of "cones under  $d$ ",  $\mathcal{C}_{d/}$ .

**Definition:** A **Colimit** of  $d$  is an initial object of  $\mathcal{C}_{d/}$ . In particular, by definition the collection of candidates for a colimit form a contractible space.

Similar for **Limits**.

# Quasi-categories as $\infty$ -categories: Interpreting Homotopy Colimits.

**Prop:**  $F : J \rightarrow E$  simplicial functor between simplicial categories enriched in Kan complexes. Let  $C$  be an object of  $E$  together with a compatible family of maps  $\{\eta_j : F(j) \rightarrow C\}_{j \in J}$ . Then

$C$  is a homotopy colimit of  $F$  iff the induced map of simplicial sets  $N_{\Delta}(J)^{\triangleright} \rightarrow N_{\Delta}(E)$  is a colimit diagram.

$$\begin{array}{ccc}
 \text{Example: } \text{Coeq}^{\mathcal{S}}\left(* \begin{array}{c} \xrightarrow{Id} \\ \xrightarrow{Id} \end{array} *\right) = \text{colim}^{\mathcal{S}} & \begin{array}{c} * \amalg * \longrightarrow * \\ \downarrow \\ * \end{array} & = \\
 \\
 = \text{hcolim.} & \begin{array}{c} * \amalg * \longrightarrow * \\ \downarrow \\ * \end{array} & = \text{colim.} & \begin{array}{c} * \amalg * \longrightarrow \Delta^1 \\ \downarrow \\ \Delta^1 \end{array} & = S^1
 \end{array}$$

**Example:**  $x : * \rightarrow X$  in  $\mathcal{S}$ .  $* \times_X^h * \simeq \Omega_x X$

# Quasi-categories as $\infty$ -categories: Dwyer-Kan Localization

**Now:** We have  $\text{Cat}_\infty$  and can talk about limits and colimits inside.

A usual category  $\rightsquigarrow N(A)$   $\infty$ -category with unique compositions.

**Definition:**  $W$  class of morphisms in  $A$ . An  $\infty$ -localization of  $A$  along  $W$ , is a quasi-category  $N(A)[W^{-1}]_\infty$  + a map of simplicial sets  $\ell : N(A) \rightarrow N(A)[W^{-1}]_\infty$  such that for any quasi-category  $\mathcal{C}$ ,

$$\text{Fun}(N(A)[W^{-1}]_\infty, \mathcal{C}) \rightarrow \text{Fun}(N(A), \mathcal{C})$$

is fully faithful with essential image =  $\infty$ -functors sending  $W$  to equivalences in  $\mathcal{C}$ .

**Remark:**  $N(A)[W^{-1}]_\infty$  can be obtained as a pushout in  $\text{Cat}_\infty$ .

$$\begin{array}{ccc} \coprod_{w \in W} N([1]) & \longrightarrow & N(A) \\ \downarrow & & \downarrow \ell \\ \coprod_{w \in W} N(J) & \longrightarrow & N(A)[W^{-1}]_\infty \end{array} \quad J = \text{groupoid 2 objects, 1 iso.}$$

# Quasi-categories as $\infty$ -categories: Dwyer-Kan Localization

**Quillen, Dwyer-Kan:**  $\mathcal{M}$  simplicial model category with  $W$  weak-equivalences. There there is a chain of equivalences of  $\infty$ -categories

$$N(\mathcal{M})[W^{-1}]_{\infty} \underset{\substack{\simeq \\ \text{cof. repl.}}}{\square} N(\mathcal{M}^c)[W_c^{-1}]_{\infty} \simeq N_{\Delta}(E_{\mathcal{M}})$$

**Definition:** We use the terminology **underlying  $\infty$ -category of  $\mathcal{M}$**  to address one of these equivalent quasi-categories.

**Remark:**  $h(N(\mathcal{M})[W^{-1}]_{\infty})$  is the usual homotopy category of  $\mathcal{M}$  (Gabriel-Zisman localization).

## Examples

- $\infty$ -derived category of a ring  $R$ ,  $D_{\infty}(R) := N(\mathcal{M})[W^{-1}]_{\infty}$  with  $\mathcal{M}$ = strict category of chain complexes of  $R$ -modules.  $W$ =quasi-iso's.
- $\infty$ -category of cdga's/ $\mathbb{C}$ ,  $\text{cdga}_{\mathbb{C}}^{\infty} := N(\mathcal{M})[W^{-1}]_{\infty}$  with  $\mathcal{M}$ = strict cdga's,  $W$ =quasi-iso's.

# Yoneda and Presheaves

**Definition:** quasi-category of presheaves of spaces on quasi-category  $\mathcal{C}$

$$\mathcal{P}(\mathcal{C}) := \text{Fun}(\mathcal{C}^{op}, \mathcal{S})$$

**Seen before:** A usual category,  $\mathcal{C} := \mathbf{N}(A)$  then  $\mathcal{P}(\mathcal{C}) \simeq \mathbf{N}_{\Delta}(E_{\mathcal{M}})$ ,  $\mathcal{M} :=$  simplicial presheaves on  $A$ .

**Prop (Yoneda):** For any quasi-category  $\mathcal{C}$  there exists a fully faithful  $\infty$ -functor

$$j : \mathcal{C} \rightarrow \mathcal{P}(\mathcal{C})$$

with the following **universal property:** If  $\mathcal{D}$  has all colimits, then the composition

$$\underbrace{\text{Fun}^L}_{\text{Left-adjoints}}(\mathcal{P}(\mathcal{C}), \mathcal{D}) \rightarrow \text{Fun}(\mathcal{C}, \mathcal{D})$$

is an equivalence of  $\infty$ -categories

**Remark:** To construct  $j$  one needs to exhibit a cocartesian fibration  $\mathcal{N} \rightarrow \mathcal{C}^{op} \times \mathcal{C}$ .  $\mathcal{N}$  is given by the  $\infty$ -category of *Twisted arrows* in  $\mathcal{C}$ .

# How to construct Yoneda?

**More generally:** How to construct  $\infty$ -functors? Not enough to say what happens on objects and 1-morphisms. Need to explain the action on  $n$ -simplexes.

**Idea:** Re-organize an  $\infty$ -functor  $F : \mathcal{C} \rightarrow \mathcal{D}$  as a "family with a connection data".

$$\begin{array}{ccc} & & \\ & \downarrow & \downarrow \\ \mathcal{N} & \mathcal{C} = p^{-1}(0) & \mathcal{D} = p^{-1}(1) \\ \downarrow p & & \\ \Delta^1 & 0 \longrightarrow & 1 \end{array}$$

with  $\text{Map}_{\mathcal{N}}((0, c), (1, d)) \simeq \text{Map}_{\mathcal{D}}(F(c), d)$

**Key point:** Not all families over  $\Delta^1$  define  $\infty$ -functors. Need connection data to jump between fibers.

**coCartesian fibration:** = family + appropriate connection data.

# Grothendieck construction

**Previous example:** Single  $\infty$ -functor  $\leftrightarrow$  Family over  $\Delta^1$ . This dictionary can be explained by an  $\infty$ -categories

$$\mathrm{coCart}/\Delta^1 \simeq \mathrm{Fun}(\Delta^1, \mathrm{Cat}_\infty)$$

**More generally:** Replacing  $\Delta^1$  by any simplicial set  $S$ , it is still true that diagrams of  $\infty$ -categories indexed by  $S$  can be exhibited as families over  $S$  with connection:

$$\mathrm{coCart}/S \simeq \mathrm{Fun}(S, \mathrm{Cat}_\infty)$$

The way the connection data is implemented in the definition of cocartesian fibration ensures the functoriality up to homotopy.

**Back to Yoneda:** Need to construct  $\infty$ -functor  $\mathcal{C} \times \mathcal{C}^{op} \rightarrow \mathcal{S} \subseteq \mathrm{Cat}_\infty$  sending  $(X, Y) \rightarrow \mathrm{Map}_{\mathcal{C}}(X, Y)$ . Instead, construct the appropriate family over  $\mathcal{C} \times \mathcal{C}^{op}$ .

**Exercise:** Formulate the notion of [pair of adjoint functors](#) as a family over  $\Delta^1$  where the connection data works in both directions-[bicartesian fibration](#).



# Presentable Categories

$\kappa$  strongly inaccessible card.  $\leftrightarrow \mathcal{U} :=$  sets of rank  $< \kappa$ , universe

$\mathcal{U} \subseteq \mathcal{V}$  universes  $\rightsquigarrow \mathcal{U}$ -small SSets  $\leftrightarrow \mathcal{V}$ -small SSets  $:= \mathcal{U} - \text{big}$

$$\text{Cat}_{\infty}^{\text{small}} \leftrightarrow \text{Cat}_{\infty}^{\text{big}}$$

**Presentable data:**  $\mathcal{U}$ -Big Data determined by  $\mathcal{U}$ -Small Data.

**Definition:**  $\mathcal{C}_0$  small  $\infty$ -category.  $\kappa' < \kappa$ .  $\mathcal{C}_0$  is  $\kappa'$ -filtered if for every  $\kappa'$ -small simplicial set  $K$  and diagram  $F : K \rightarrow \mathcal{C}_0$  there exists an extension  $K^{\triangleright} \rightarrow \mathcal{C}_0$ .

**Definition:**  $\mathcal{C}_0$  small  $\infty$ -category.  $\kappa' < \kappa$ .  $\text{Ind}_{\kappa'}(\mathcal{C}_0) :=$  full subcategory of  $\mathcal{P}(\mathcal{C}_0)$  containing representables + stable under  $\kappa'$ -filtered colimits.

**Definition:**  $\mathcal{C}$   $\infty$ -category is **presentable** if  $\mathcal{C}$  has all small colimits and is of the form  $\text{Ind}_{\kappa'}(\mathcal{C}_0)$  for some small  $\infty$ -category  $\mathcal{C}_0$ .

**Examples:**  $\mathcal{S}$  is presentable, of the form  $\text{Ind}_{\omega}$  of (homotopy) finite CW-complexes. All categories of presheaves are presentable.

# Presentable Categories: Adjoint Functor Theorem

**Prop. (Adjoint Functor Theorem).**  $F : \mathcal{C} \rightarrow \mathcal{D}$   $\infty$ - functor between presentable  $\infty$ -categories. Suppose  $F$  commutes with all small colimits. Then  $F$  admits a right adjoint.

**Definition:**  $\mathbf{Pr}^{\mathbf{L}}$  is the (non-full) subcategory of  $\mathbf{Cat}_{\infty}^{\mathbf{big}}$  spanned by presentable  $\infty$ -categories and colimit preserving functors.

**Remark:** One can also define a (non-full) subcategory  $\mathbf{Pr}^{\mathbf{R}} \subseteq \mathbf{Cat}_{\infty}^{\mathbf{big}}$  spanned by presentable categories and functors which preserve limits and  $\kappa$ 's filtered colimits for some cardinal  $\kappa' < \kappa$ . We have  $(\mathbf{Pr}^{\mathbf{R}})^{op} \simeq \mathbf{Pr}^{\mathbf{L}}$

**Prop:** The inclusions  $\mathbf{Pr}^{\mathbf{R}}, \mathbf{Pr}^{\mathbf{L}} \subseteq \mathbf{Cat}_{\infty}^{\mathbf{big}}$  preserve all limits.

**Prop:**  $\mathcal{C}$  and  $\mathcal{D}$  presentable. Then  $\mathbf{Fun}^{\mathbf{L}}(\mathcal{C}, \mathcal{D})$  is presentable.

# Presentable Categories: Localizations

**Prop. (Presentable Localizations):**  $\mathcal{C}$  presentable  $\infty$ -category,  $W$  small collection of 1-morphisms in  $\mathcal{C}$ .  $\mathcal{C}^{W\text{-local}} :=$  full subcategory of  $\mathcal{C}$  spanned by  $X$  such that for every  $S \rightarrow S'$  in  $W$  the map

$$\mathrm{Map}_{\mathcal{C}}(S', X) \rightarrow \mathrm{Map}_{\mathcal{C}}(S, X)$$

is an equivalence. Then  $\mathcal{C}^{W\text{-local}}$  is presentable and the inclusion  $\mathcal{C}^{W\text{-local}} \subseteq \mathcal{C}$  admits a left adjoint  $\mathcal{C} \rightarrow \mathcal{C}^{W\text{-local}}$  which exhibits  $\mathcal{C}^{W\text{-local}}$  as an  $\infty$ -localization of  $\mathcal{C}$  along  $W$  **internal** to the theory of presentable  $\infty$ -categories.

**Examples:** Produce examples of presentable categories as presentable localizations of presheaves. In fact:

**Prop:**  $\mathcal{C}$  is presentable iff it is a presentable localization of some  $\mathcal{P}(\mathcal{C}_0)$ .

**Prop:** A usual category. Presentable localizations of  $\mathcal{P}(N(A))$  correspond via the rectification result to Bousfield localizations of simplicial presheaves on  $A$ .

**What is a symmetric monoidal structure?:** A usual category + unit  $*$   $\rightarrow A$  +  $n$ -ary operations  $\otimes^n : \underbrace{A \times \dots \times A}_n \rightarrow A$  + associativity constrains + symmetry constrains under the action of the  $\Sigma_n$ .

**Segal:** This combinatorial pattern is described by the category of finite pointed sets  $\text{Fin}_*$  and the data of a symmetric monoidal category is the data of a pseudo-functor

$$A^\otimes : \text{Fin}_* \rightarrow \text{Cat}, \text{ with property } A^\otimes(\{n\}_0) = \underbrace{A \times \dots \times A}_n$$

**Example:**  $\{0, 1, 2\} \rightarrow \{0, 1\}$   $1, 2 \mapsto 1, 0 \mapsto 0$ , Image under  $A^\otimes$  gives operation  $A \times A \rightarrow A$ . Image of permutation of 1 by 2,  $\{0, 1, 2\} \rightarrow \{0, 1, 2\}$  gives symmetry.

**Symmetric monoidal  $\infty$ -category :=  $\infty$ -functor**

$$c^\otimes : \mathcal{N}(\text{Fin}_*) \rightarrow \text{Cat}_\infty, \text{ with } c^\otimes(\{n\}_0) \rightarrow \prod_{i=1}^n c^\otimes(\{0, 1\}) \text{ eq.}$$

## Some subtle points:

- **"The diagram commutes"**. The claim does not make sense unless the  $n$ -simplex providing the commutativity is exhibited. One should look for an explicit construction of the simplex or for a universal property or "machine" defining it.
- **"The  $\infty$ -functor  $F : \mathcal{C} \rightarrow \mathcal{D}$  doing this on objects and that on 1-morphisms"**. Such a description should always be understood as informal. The only construction with mathematical sense is that of a map of simplicial sets  $F : \mathcal{C} \rightarrow \mathcal{D}$  which specifies what happens with all simplexes. Typically we construct  $\infty$ -functors via some universal property involved, via Quillen adjunctions or via the exhibition of a cartesian/cocartesian fibration.
- **"The object has an (associative, commutative, Lie)-algebra structure"**. The exhibition and construction of such structures is in general very subtle.

# Main $\infty$ -categories in these lectures:

- Spaces  $\mathcal{S}$ ,  $\infty$ -categories  $\text{Cat}_{\infty}^{\text{small}}$ ,  $\text{Cat}_{\infty}^{\text{big}}$ , presentable  $\infty$ -categories  $\text{Pr}^{\text{L}}$ , stable presentable  $\text{Pr}_{\text{Stb}}^{\text{L}}$
- cdga's  $\text{cdga}_{\mathbb{C}}^{\infty}$ , derived schemes/manifolds/stacks/prestacks;
- Spectra  $\text{Sp}$ , derived category of a ring  $R$ ,  $D_{\infty}(R)$ ,  $\text{Qcoh}_{\infty}(X)$ ,  $\text{Perf}(X)$  (quasi-coherent and perfect complexes),
- $\infty$ -categories of algebra-objects and modules over them in all the  $\infty$ -categories above.
- $(\infty, n)$ -categories of correspondences, Lagrangian correspondences, cobordisms, etc.

Ready for action!